

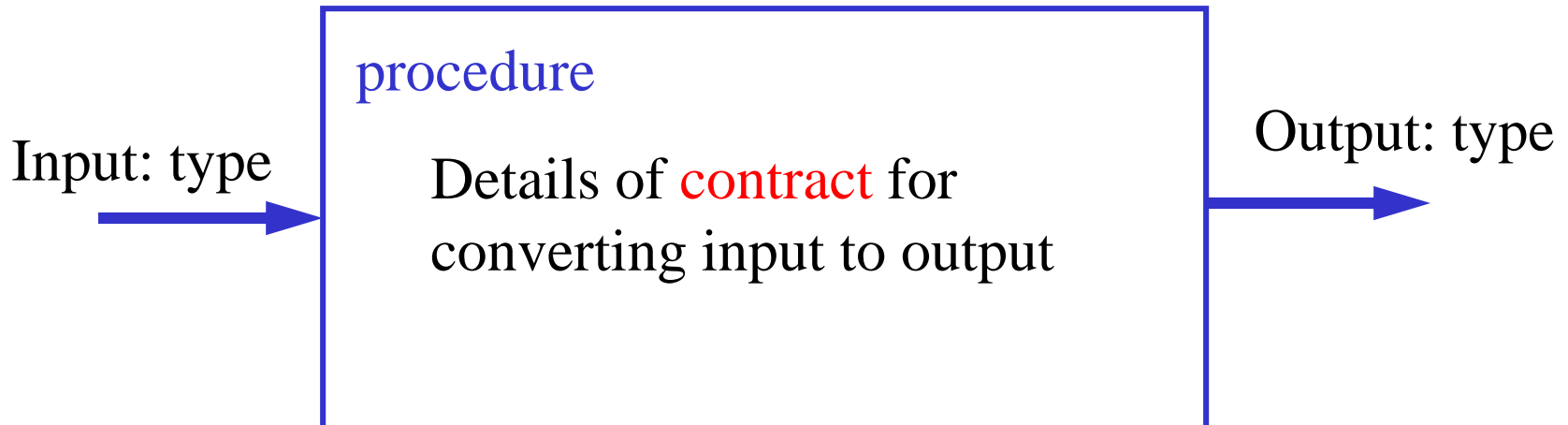
# Higher-Order Procedures

- Today's topics
  - Procedural abstractions
  - Capturing patterns across procedures – Higher Order Procedures

# Procedural abstraction

- **Process of procedural abstraction**

- Define formal parameters, capture pattern of computation as a process in body of procedure
- Give procedure a name
- Hide implementation details from user, who just invokes name to apply procedure



# Procedural abstraction example: sqrt

To find an approximation of square root of  $x$ :

- Make a guess  $G$
- Improve the guess by averaging  $G$  and  $x/G$
- Keep improving the guess until it is good enough

```
(define try (lambda (guess x)
              (if (good-enuf? guess x)
                  guess
                  (try (improve guess x) x))))

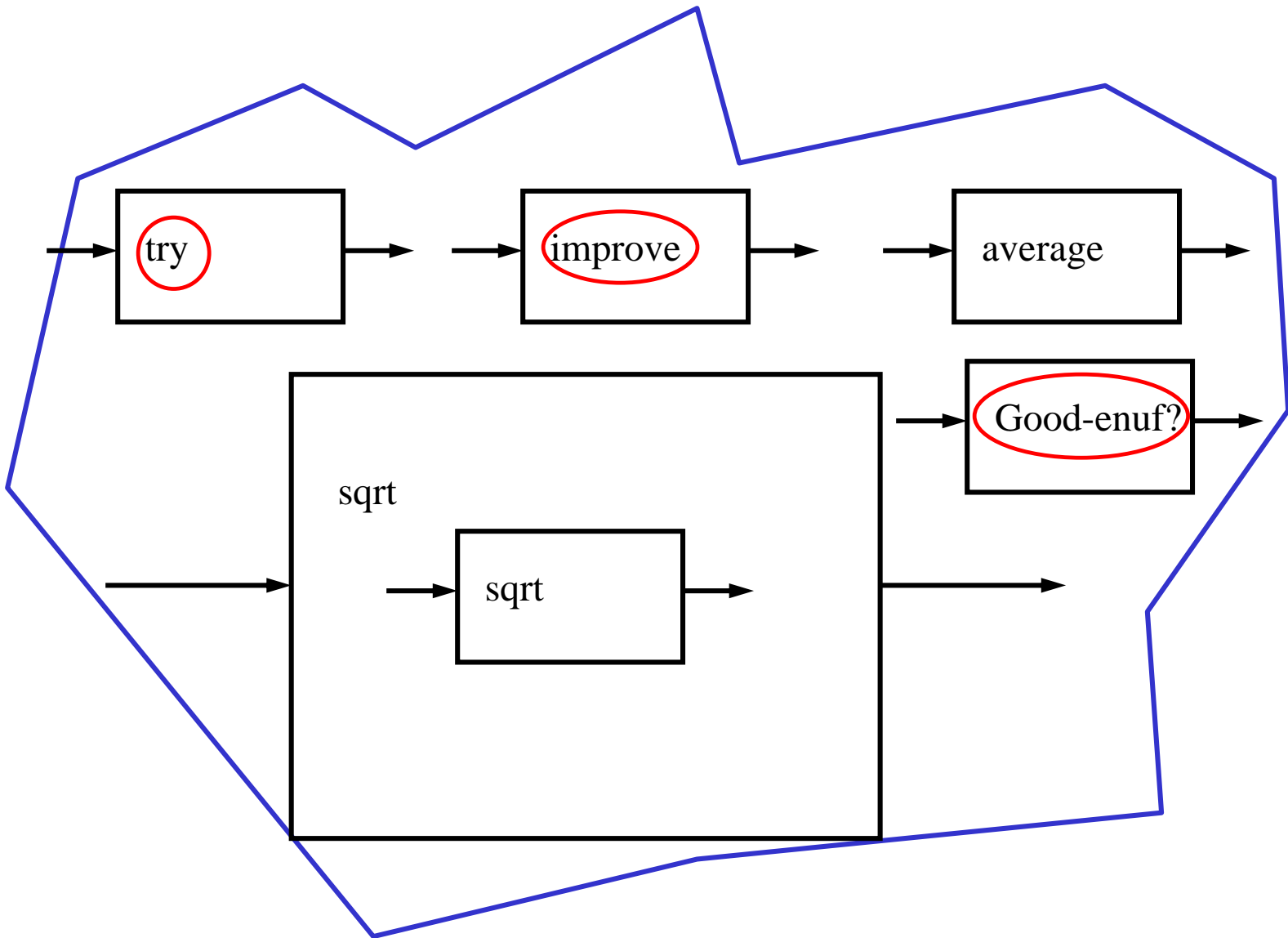
(define good-enuf? (lambda (guess x)
                     (< (abs (- (square guess) x)) 0.001)))

(define improve (lambda (guess x)
                  (average guess (/ x guess))))

(define average (lambda (a b) (/ (+ a b) 2)))

(define sqrt (lambda (x) (try 1 x)))
```

# The universe of procedures for sqrt



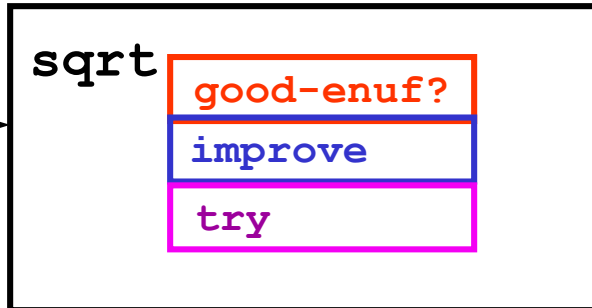
# sqrt - Block Structure

```
(define sqrt
```

```
(lambda (x)
  (define good-enuf?
    (lambda (guess)
      (< (abs (- (square guess) x))
        0.001)))
  (define improve
    (lambda (guess)
      (average guess (/ x guess))))
  (define try
    (lambda (guess)
      (if (good-enuf? guess)
          guess
          (try (improve guess)))))
  (try 1))
```

```
)
```

**x: number** →



→ **√x: number**

# Typecasting

- We are going to find that it is convenient to reason about procedures (and data structures) in terms of the number and kinds of arguments, and the kind of output they produce
- We call this **typing** of a procedure or data structure

# Types – a motivation

```
(+ 5 10) ==> 15
```

```
(+ "hi" 5)
```

**;The object "hi", passed as the first argument to integer-add, is not the correct type**

- Addition is not defined for strings

# Types – simple data

- We want to collect a taxonomy of expression types:
  - Simple Data
    - Number
      - Integer
      - Real
      - Rational
    - String
    - Boolean
    - Names (symbols)
- We will use this for notational purposes, to reason about our code. Scheme checks types of arguments for built-in procedures, but *not for user-defined ones*.



# Types – procedures

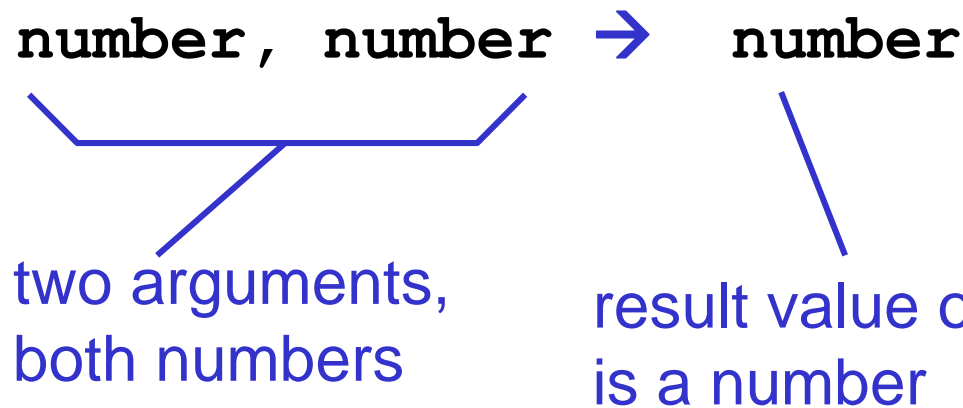
- Because procedures operate on objects and return values, we can define their types as well.
- We will denote a procedures type by indicating the types of each of its arguments, and the type of the returned value, plus the symbol  $\rightarrow$  to indicate that the arguments are mapped to the return value
- E.g. **number  $\rightarrow$  number** specifies a procedure that takes a number as input, and returns a number as value

# Types

- `(+ 5 10) ==> 15`  
`(+ "hi" 5)`

**;The object "hi", passed as the first argument to integer-add, is not the correct type**

- Addition is not defined for strings
- The **type** of the integer-add procedure is





# Types, precisely

- A type describes a **set** of scheme **values**

- **number** → **number** describes the set:

all procedures, whose result is a number,  
which require one argument that must be a number

- **Every** scheme value has a type

- Some values can be described by multiple types
- If so, choose the type which describes the largest set

- Special form keywords like **define** do not name values

- therefore special form keywords **have no type**

# Your turn

- The following expressions evaluate to values of what type?

```
(lambda (a b c) (if (> a 0) (+ b c) (- b c)))
```

number, number, number → number

```
(lambda (p) (if p "hi" "bye"))
```

Boolean → string

```
(* 3.14 (* 2 5))
```

number

# Summary of types

- type: a set of values
- every value has a type
- procedure types (types which include  $\rightarrow$ ) indicate
  - number of arguments required
  - type of each argument
  - type of result of the procedure
- Types: a mathematical theory for reasoning **efficiently** about programs
  - useful for preventing certain common types of errors
  - basis for many analysis and optimization algorithms

# What is procedure abstraction?

Capture a common pattern

(`* 2 2`)

(`* 57 57`)

(`* k k`)

(`lambda (x) (* x x)`)



Formal parameter for pattern



Actual pattern

Give it a name (`define square (lambda (x) (* x x))`)

Note the type: `number`  $\rightarrow$  `number`

# Other common patterns

- $1 + 2 + \dots + 100$
- $1 + 4 + 9 + \dots + 100^2$
- $1 + 1/3^2 + 1/5^2 + \dots + 1/101^2 (= \pi^2/8)$

$$\sum_{k=1}^{100} k$$
$$\sum_{k=1}^{100} k^2$$
$$\sum_{k=1, \text{odd}}^{101} k^{-2}$$

```
(define (sum-integers a b)
  (if (> a b)
      0
      (+ a (sum-integers (+ 1 a) b))))
```

```
(define (sum-squares a b)
  (if (> a b)
      0
      (+ (square a)
         (sum-squares (+ 1 a) b))))
```

```
(define (pi-sum a b)
  (if (> a b)
      0
      (+ (/ 1 (square a))
         (pi-sum (+ a 2) b))))
```

```
(define (sum term a next b)
```

```
  (if (> a b)
```

```
      0
```

```
      (+ (term a)
```

```
         (sum term (next a) next b))))
```



# Let's examine this new procedure

```
(define (sum term a next b)  
  (if (> a b)  
      0  
      (+ (term a)  
          (sum term (next a) next b))))
```

What is the type of this procedure?

$(\text{num} \rightarrow \text{num}, \text{num}, \text{num} \rightarrow \text{num}, \text{num}) \rightarrow \text{num}$

1. What type is the output?
2. How many arguments?
3. What type is each argument?

# Higher order procedures

- A higher order procedure:  
takes a procedure as an *argument* or returns one as a *value*

```
(define (sum-integers a b)
  (if (> a b)
      0
      (+ a (sum-integers (+ 1 a) b))))
```

```
(define (sum term a next b)
  (if (> a b)
      0
      (+ (term a) (sum term (next a) next b))))
```

```
(define (sum-integers1 a b)
  (sum (lambda (x) x) a (lambda (x) (+ x 1)) b))
```

# Higher order procedures

```
(define (sum-squares a b)
  (if (> a b)
      0
      (+ (square a)
          (sum-squares (+ 1 a) b))))
```

```
(define (sum term a next b)
  (if (> a b)
      0
      (+ (term a) (sum term (next a) next b))))
```

```
(define (sum-squares1 a b)
  (sum square a (lambda (x) (+ x 1)) b))
```

# Higher order procedures

```
(define (pi-sum a b)
  (if (> a b)
      0
      (+ (/ 1 (square a))
          (pi-sum (+ a 2) b))))
```

```
(define (sum term a next b)
  (if (> a b)
      0
      (+ (term a) (sum term (next a) next b))))
```

```
(define (pi-sum1 a b)
  (sum (lambda (x) (/ 1 (square x))) a
        (lambda (x) (+ x 2)) b))
```

# Higher order procedures

- Takes a procedure as an argument or returns one as a value

```
(define (sum-integers1 a b)
  (sum (lambda (x) x) a (lambda (x) (+ x 1)) b))
```

```
(define (sum-squares1 a b)
  (sum square a (lambda (x) (+ x 1)) b))
```

```
(define (add1 x) (+ x 1))
```

```
(define (sum-squares1 a b) (sum square a add1 b))
```

```
(define (pi-sum1 a b)
  (sum (lambda (x) (/ 1 (square x))) a
  (lambda (x) (+ x 2)) b))
```

```
(define (add2 x) (+ x 2))
```

```
(define (pi-sum1 a b)
  (sum (lambda (x) (/ 1 (square x))) a add2 b))
```

# Returning A Procedure As A Value

```
(define (add1 x) (+ x 1))
```

```
(define (add2 x) (+ x 2))
```

```
(define incrementby (lambda (n) . . . ))
```

```
(define add1 (incrementby 1))
```

```
(define add2 (incrementby 2))
```

```
. . .
```

```
(define add37.5 (incrementby 37.5))
```

```
incrementby: # → (# → #)
```

# Returning A Procedure As A Value

```
(define incrementby  
  (lambda (n) [REDACTED]))
```

```
(incrementby  
 (lambda (n) (lambda (x) (+ x n)))) 2) →
```

```
(lambda (x) (+ x 2))
```

`(incrementby 2)` → a *procedure* of one var (`x`) that increments `x` by 2

`((incrementby 3) 4)` → ?

`( (lambda (x) (+ x 3)) 4)` →

# Nano-Quiz/Lecture Problem

```
(define incrementby  
  (lambda (n) (lambda (x) (+ x n))))
```

```
(define f1 (incrementby 6)) → ?
```

```
(f1 4) →
```

```
(define f2 (lambda (x) (incrementby 6))) → ?
```

```
(f2 4) → ?
```

```
((f2 4) 6) → ?
```



# Procedures as values: Derivatives

$$f : x \rightarrow x^2$$

$$f : x \rightarrow x^3$$

$$f' : x \rightarrow 2x$$

$$f' : x \rightarrow 3x^2$$

- Taking the derivative is a function:  $D(f) = f'$
- What is its *type*?

$$D : (\# \rightarrow \#) \rightarrow (\# \rightarrow \#)$$

# Computing derivatives

- A good approximation:

$$Df(x) \approx \frac{\underline{f(x + \epsilon)} - \underline{f(x)}}{\underline{\epsilon}}$$

```
(define deriv
  (lambda (f)
    { (lambda (x) (/ (- (f (+ x epsilon)) (f x))
                     epsilon)) )
```

(number → number) → (number → number)



# Finding fixed points of functions

Square root of  $x$  is defined by  $\sqrt{x} = x / \sqrt{x}$

Think of as a transformation  $f : y \rightarrow \frac{x}{y}$  then if we can find a  $y = \sqrt{x}$ , then  $f(y) = y$ , and such a  $y$  is called a fixed point of  $f$ .

- Here's a common way of finding fixed points

- Given a guess  $x_1$ , let new guess be  $f(x_1)$
- Keep computing  $f$  of last guess, till close enough

```
(define (close? u v) (< (abs (- u v)) 0.0001))
```

```
(define (fixed-point f i-guess)
```

```
  (define (try g)
```

```
    (if (close? (f g) g)
```

```
        (f g)
```

```
        (try (f g))))
```

```
(try i-guess))
```

# Using fixed points

```
(fixed-point (lambda (x) (+ 1 (/ 1 x))) 1)  
→ 1.6180
```

or  $x = 1 + 1/x$  when  $x = (1 + \sqrt{5})/2$

```
(define (sqrt x)  
  (fixed-point  
    (lambda (y) (/ x y))  
    1))
```

$$y = \frac{x}{y}$$

$$y^2 = x$$

$$y = \sqrt{x}$$

Unfortunately if we try (sqrt 2), this oscillates between 1, 2, 1, 2,

```
(define (fixed-point f i-guess)  
  (define (try g)  
    (if (close? (f g) g)  
        (f g)  
        (try (f g))))  
  (try i-guess))
```

# So damp out the oscillation

```
(define (average-damp f)
  (lambda (x)
    (average x (f x))))
```

Check out the type:

$(\text{number} \rightarrow \text{number}) \rightarrow (\text{number} \rightarrow \text{number})$

that is, this takes a procedure as input, and returns a **NEW** procedure as output!!!

- ((average-damp square) 10)
- ((lambda (x) (average x (square x))) 10)
- (average 10 (square 10))
- 55

**... which gives us a clean version of sqrt**

```
(define (sqrt x)
  (fixed-point
    (average-damp
      (lambda (y) (/ x y))))
  1))
```

Compare this to Heron's algorithm (the one we saw earlier)  
– same process, but ideas intertwined with code

```
(define (cbrt x)
  (fixed-point
    (average-damp
      (lambda (y) (/ x (square y)))))
  1))
```

# Procedures as arguments: a more complex example

- ```
(define compose (lambda (f g x) (f (g x))))  
  (compose square double 3)  
  (square (double 3))  
  (square (* 3 2))  
  (square 6)  
  (* 6 6)  
  36
```

What is the type of `compose`? Is it:

`(number → number), (number → number), number → number`

**No! Nothing in `compose` requires a number**



# Compose works on other types too

```
(define compose (lambda (f g x) (f (g x))))
```

```
(compose  
  (lambda (p) (if p "hi" "bye"))  
  (lambda (x) (> x 0))  
  -5  
) ==> "bye"
```

boolean → string  
number → boolean  
number  
result: a string

Will any call to compose work?

```
(compose < square 5)
```

*wrong number of args to <*

*<: number, number → boolean*

```
(compose square double "hi")
```

*wrong type of arg to double*

*double: number → number*

# Type of compose

```
(define compose (lambda (f g x) (f (g x))))
```

- Use **type variables**.

**compose:**     (B → C), (A → B), A → C

- Meaning of type variables:

All places where a given type variable appears must match when you fill in the actual operand types

- The constraints are:

- F and G must be functions of one argument
- the argument type of G matches the type of X
- the argument type of F matches the result type of G
- the result type of compose is the result type of F

# Higher order procedures

- Procedures may be passed in as arguments
- Procedures may be returned as values
- Procedures may be used as parts of data structures
- Procedures are first class objects in Scheme!!